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Abstract

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MATHEMATICS

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ON NONTRIVIAL CONSTRUCTIVE MAPPINGS OF CERTAIN SETS

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The present note belongs to the constructive trend in mathematics ⁽¹⁾. The term **set** is used here in the same sense as in papers ^(2,3). We shall assume that, for the elements of each of the sets considered here, a reflexive, symmetric, and transitive equality relation has been introduced (by means of a two-place predicate), and elements that stand in the equality relation will be called **equal**. If $\mathfrak{M}$ is a set in an alphabet A , then A will be called the **alphabet of the set** $\mathfrak{M}$. Let $\mathfrak{M}$ and $\mathfrak{N}$ be two sets. By a **partial mapping** of the set $\mathfrak{M}$ into the set $\mathfrak{N}$ we shall mean any normal algorithm ⁽⁴⁾ over the union of the alphabets of these sets that transforms those elements of the set $\mathfrak{M}$ to which it is applicable into elements of the set $\mathfrak{N}$. A partial mapping $\mathfrak{A}$ of the set $\mathfrak{M}$ into the set $\mathfrak{N}$ will be called a **complete mapping** of $\mathfrak{M}$ into $\mathfrak{N}$ if the algorithm $\mathfrak{A}$ is applicable to every element of the set $\mathfrak{M}$. A partial mapping $\mathfrak{A}$ of the set $\mathfrak{M}$ into the set $\mathfrak{N}$ will be called an **equality-preserving mapping** if, for any elements P and Q of the set $\mathfrak{M}$, from the facts that P and Q are equal elements and that the algorithm $\mathfrak{A}$ is applicable both to P and to Q , it follows that the elements $\mathfrak{A}(P)$ and $\mathfrak{A}(Q)$ of the set $\mathfrak{N}$ are also equal. A partial mapping $\mathfrak{A}$ of the set $\mathfrak{M}$ into the set $\mathfrak{N}$ will be called **nontrivial** if it is false that there do not exist elements P and Q of the set $\mathfrak{M}$ such that the algorithm $\mathfrak{A}$ is applicable both to P and to Q and the elements $\mathfrak{A}(P)$ and $\mathfrak{A}(Q)$ of the set $\mathfrak{N}$ are not equal. A **trivial** partial mapping will be one that is not nontrivial.

There are known examples of such sets $\mathfrak{M}, \mathfrak{N}$ that every complete mapping of $\mathfrak{M}$ into $\mathfrak{N}$ preserving equality transforms all elements of the set $\mathfrak{M}$ into elements of the set $\mathfrak{N}$ equal to one another, i.e. every complete mapping of $\mathfrak{M}$ into $\mathfrak{N}$ preserving equality is trivial. Thus, for example, every complete mapping of the set of constructive real numbers—duplexes ⁽³⁾—into the set of integers that transforms equal real numbers into equal integers is trivial. This follows from the known properties of constructive functions of a real variable (see Corollary {3} in § 4 in ⁽⁵⁾). Here we shall give other analogous examples, and moreover shall establish the impossibility not only of complete nontrivial equality-preserving mappings, but also of nontrivial partial mappings preserving equality and satisfying certain conditions. These examples are based on the

result formulated in Theorem 1.

Let $\mathfrak{M}$ be a set and P an element of this set. A normal algorithm $\mathfrak{A}$ over the alphabet of this set will be called an **algorithm complete on P in $\mathfrak{M}$** if it is applicable to every element of the set $\mathfrak{M}$ equal to the element P . In the same sense we shall use the term **partial mapping complete on P in $\mathfrak{M}$** . A normal algorithm $\mathfrak{A}$ over the alphabet of the set $\mathfrak{M}$ will be called an **algorithm empty on P in $\mathfrak{M}$** if it is not applicable to any element of the set $\mathfrak{M}$ equal to the element P . An element P of the set $\mathfrak{M}$ will be called a **node** if, for no element Q of this set, there exists a normal algorithm complete on P in $\mathfrak{M}$ and empty on Q .

in $\mathfrak{M}$. We shall say that a set $\mathfrak{M}$ has an **enumerable equality condition** if there exists a normal algorithm applicable to a pair of elements P, Q of the set $\mathfrak{M}$ if and only if the elements P and Q are equal. We shall say that a set $\mathfrak{M}$ has an **enumerable inequality condition** if there exists a normal algorithm applicable to a pair of elements P, Q of the set $\mathfrak{M}$ if and only if the elements P and Q are not equal.

Theorem 1. *Let $\mathfrak{M}$ be a set and P a nodal element of this set. Whatever the set $\mathfrak{N}$, having an enumerable equality condition or an enumerable inequality condition, it is impossible to construct a nontrivial partial mapping of $\mathfrak{M}$ into $\mathfrak{N}$ that preserves equality and is complete on P in $\mathfrak{M}$.*

We shall use the term **quasnumber** in the same sense as in [5]. Let $\mathfrak{M}$ be a set. We shall say that $\mathfrak{M}$ is an **invariant set of quasnumbers** if the following conditions hold: 1) every element of the set $\mathfrak{M}$ is a quasnumber; 2) for any quasnumbers x and y , from the fact that x and y are conditionally equal and x belongs to $\mathfrak{M}$, it follows that y also belongs to $\mathfrak{M}$. The elements of an invariant set of quasnumbers that are conditionally equal quasnumbers will be called **equal elements**.

Theorem 2. *Whatever the invariant set of quasnumbers, every quasnumber belonging to this set is a nodal element of this set.*

The set of constructive real numbers (whose equal elements are taken to be equal constructive real numbers) has an enumerable inequality condition; therefore, from Theorems 1 and 2 we obtain

Corollary 1. *Whatever the invariant set $\mathfrak{M}$ of quasnumbers, it is impossible to construct a nontrivial partial mapping of the set $\mathfrak{M}$ into the set of constructive real numbers that preserves equality and is applicable to every quasnumber conditionally equal to some quasnumber from $\mathfrak{M}$.*

By a **nondecreasing quasnumber** we shall mean any quasnumber whose basis is a nondecreasing sequence of rational numbers. Let $\mathfrak{M}$ be a set. We shall say that $\mathfrak{M}$ is an **invariant set of nondecreasing quasnumbers** if the following conditions hold: 1) every element of the set $\mathfrak{M}$ is a nondecreasing quasnumber; 2) for any nondecreasing quasnumbers x and y , from the fact

that x and y are conditionally equal and x belongs to $\mathfrak{M}$, it follows that y also belongs to $\mathfrak{M}$.

Let $\mathfrak{M}$ be an invariant set of nondecreasing quasinumbers and let x be an element of this set. We shall say that x is a **minimal element** of the set $\mathfrak{M}$ if $\mathfrak{M}$ contains no quasinumber conditionally less than x . We shall say that x is a **maximal element** of $\mathfrak{M}$ if $\mathfrak{M}$ contains no quasinumber conditionally greater than x .

Theorem 3. *The minimal element of any invariant set of nondecreasing quasinumbers is a nodal element of this set.*

Corollary 2. *Whatever the invariant set $\mathfrak{M}$ of nondecreasing quasinumbers, it is impossible to construct a nontrivial partial mapping of this set into the set of constructive real numbers that preserves equality and is applicable to all nondecreasing quasinumbers conditionally equal to some nondecreasing quasinumber from $\mathfrak{M}$ that is not a maximal element of the set $\mathfrak{M}$.*

In particular, there is no algorithm that constructs constructive limits for each of those increasing bounded-above sequences of rational numbers for which constructive limits cannot fail to exist.

Let A be an alphabet, and let α and β be distinct letters not belonging to A . Let $\mathfrak{M}$ be a set. We shall say that $\mathfrak{M}$ is a **set**

normal algorithms over the alphabet A , if every element of the set $\mathfrak{M}$ is a record of some normal algorithm in the alphabet $A \cup \{\alpha, \beta\}$.

Let $\mathfrak{M}$ be a set of normal algorithms over the alphabet A . We shall say that the elements P and Q of the set $\mathfrak{M}$ are equal if the normal algorithms $\langle P \rangle$ and $\langle Q \rangle$ are equivalent relative to A (here $\langle R \rangle$ denotes the normal algorithm in the alphabet $A \cup \{\alpha, \beta\}$ whose record is the word R). We shall say that $\mathfrak{M}$ is invariant if the following condition is satisfied: whatever the element P of the set $\mathfrak{M}$ and the normal algorithm $\mathfrak{A}$ in the alphabet $A \cup \{\alpha, \beta\}$, from the fact that $\langle P \rangle$ and $\mathfrak{A}$ are equivalent relative to A it follows that the record of the algorithm $\mathfrak{A}$ is an element of $\mathfrak{M}$.

Let $\mathfrak{A}$ and $\mathfrak{B}$ be normal algorithms over the alphabet A . We shall say that $\mathfrak{A}$ is an extension of $\mathfrak{B}$ relative to A , if for every word P in A , from the fact that $\mathfrak{B}$ is applicable to P , it follows that $\mathfrak{A}$ is applicable to P and that $\mathfrak{A}(P) = \mathfrak{B}(P)$.

Let $\mathfrak{M}$ be a set of normal algorithms over the alphabet A . We shall say that an element P of the set $\mathfrak{M}$ is minimal if for every element Q of the set $\mathfrak{M}$ the algorithm $\langle Q \rangle$ is an extension relative to A of the algorithm $\langle P \rangle$.

Theorem 4. *The minimal element of any invariant set of normal algorithms over the alphabet A is a nodal element of this set.*

Corollary 3. *Let $\mathfrak{M}$ be an invariant set of normal algorithms over the alphabet A , and let P be its minimal element. It is impossible to construct a nontrivial partial mapping of the set $\mathfrak{M}$ into the set of constructive real numbers that preserves equality and is total on P in $\mathfrak{M}$.*

Let B be an alphabet and $\mathfrak{N}$ a set. We shall say that $\mathfrak{N}$ is a set of associative calculi in the alphabet B ((⁴), Chap. 6, § 1), if every element of the set $\mathfrak{N}$ is a record of a defining system of relations of some associative calculus in B ((⁴), p. 345; henceforth, instead of the term “record of a defining system of relations of an associative calculus,” we shall use the term “record of an associative calculus”).

Let $\mathfrak{N}$ be a set of associative calculi in the alphabet B . We shall say that the elements P and Q of the set $\mathfrak{N}$ are equal if the associative calculi in B whose records are the words P and Q are isomorphic. From the results of [6] it follows that any set of associative calculi in B has an enumerable equality condition. Therefore, from Theorems 1 and 4 we obtain

Corollary 4. *Let A and B be alphabets, $\mathfrak{M}$ an invariant set of normal algorithms over the alphabet A , P the minimal element of the set $\mathfrak{M}$, $\mathfrak{N}$ a set of associative calculi in B , and $\mathfrak{A}$ a partial mapping of $\mathfrak{M}$ into $\mathfrak{N}$ that preserves equality and is total on P in $\mathfrak{M}$. If the elements Q and R of the set $\mathfrak{M}$ are such that $\mathfrak{A}(Q)$ and $\mathfrak{A}(R)$ are defined, then $\mathfrak{A}(Q)$ and $\mathfrak{A}(R)$ are equal elements of the set $\mathfrak{N}$.*

In Corollaries 3 and 4, as $\mathfrak{M}$ one may take the set of all normal algorithms over A . The minimal element of this set is the record of a normal algorithm that is not applicable to any word in A .

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